Scaling Switch-driven Flow Control with Aquarius

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- **■** Introduction
- **■** Background
- Motivation
- Aquarius Design
- Evaluation
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Introduction

Key motivations:

- End-to-end congestion controls becomes increasingly challenging to maintain effective due to the inherent feedback delay.
- Prior flow control (FC) mechanisms either lack fine-grained (i.e., per-flow granularity) control or require an impractical number of queues.

Solution:

Aquarius, a scalable solution that maintains fine-grained per-flow level control granularity with a practical number of queues.

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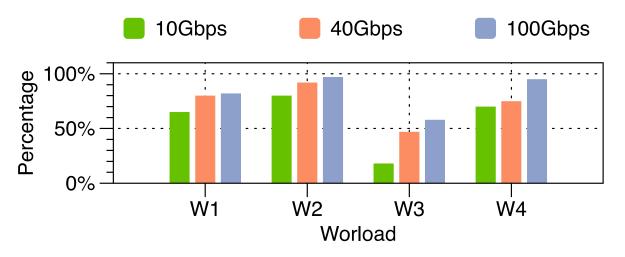
Background

Rising link speeds result in increasingly bursty traffic

Representative production datacenter workloads:

(W1) Web Server [2], (W2) Alibaba Storage [3],

(W3) Web Search [1], (W4) Facebook Hadoop [2]



Percentage of flows that finish in a single RTT (12us)

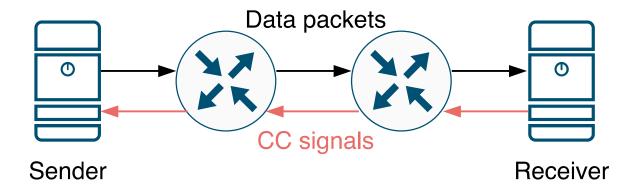
^[1] M. Alizadeh, et al. "Data center tcp (dctcp)," in Proceedings of the ACM SIGCOMM 2010 Conference.

^[2] Arjun Roy, et al. "Inside the social network's (datacenter) network", in Proceedings of the ACM SIGCOMM 2015 Conference.

^[3] Yuliang Li, et al. "HPCC: High precision congestion control", in Proceedings of the ACM SIGCOMM 2019 Conference.

Background

- End-to-end CC alone is insufficient for managing transient congestion
 - ☐ End-to-end CCs rely on receiver-echoed signals to adjust sending rates.
 - □ sender requires at least one RTT to receive feedback and loses control of flows that can complete within the first RTT.
- Per-hop flow control is necessary for handling transient congestion



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Motivation

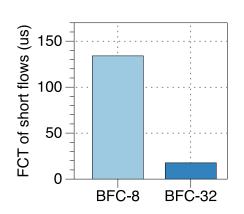
Prior flow control schemes are insufficient; they either lack fine-grained control or require an impractical number of queues.

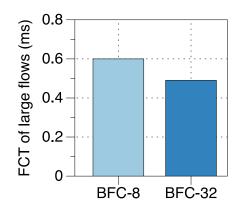
- PFC is coarse-grained
- Ideal Flow Control is fine-grained but impractical
 The ideal flow control allocates a dedicated queue to every flow, thus providing per-flow level control. However, the per-flow queue is impractical.
- Scalability issues persist in BFC [NSDI'22]

 BFC assigns a dedicated queue to each active flow if possible and enables multiple flows sharing a queue when there are no available queues.

BFC Scalability

- BFC requires more physical queues than the common switch can accommodate
 - BFC uses 32/128 queues per port.
 - (1) Majority of switches are usually equipped with 8 or fewer queues
 - (2) Physical queues are critical resources and are typically reserved for strong physical isolation between different tenants.
- BFC experiences considerable performance degradation when queues are limited



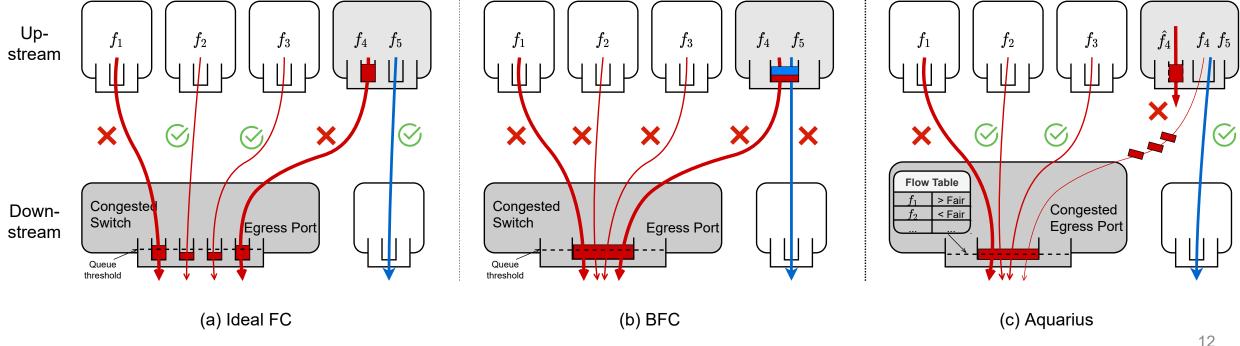


Average FCT of BFC under Web Server distribution.

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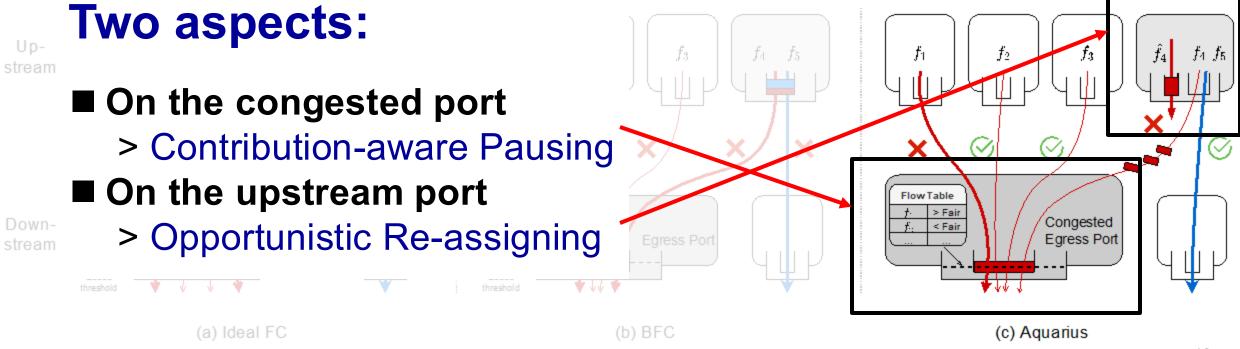
Key Idea

Approximate the ideal flow control behavior without requiring per-flow queues

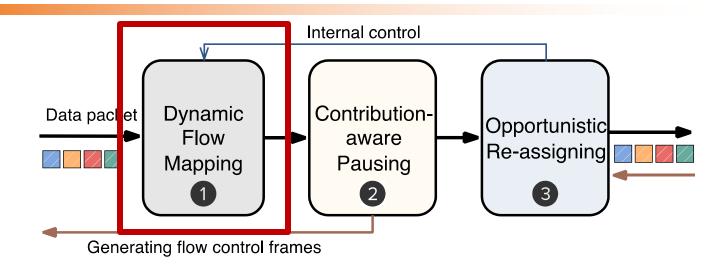


Key Idea

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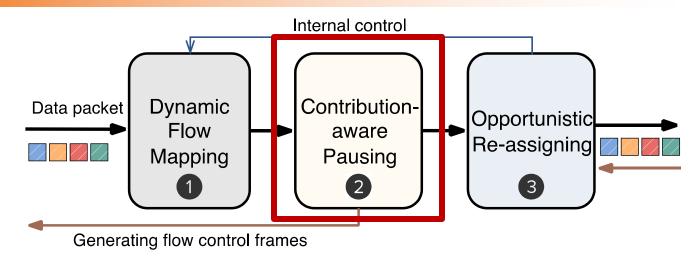


Aquarius



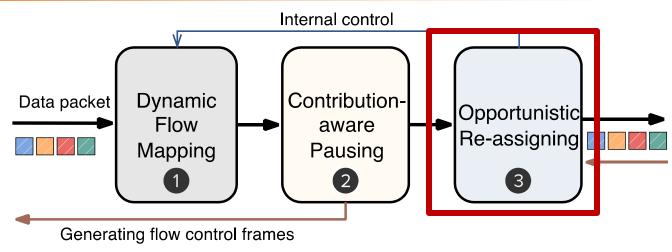
Dynamic flow mapping at every passed switch port uniformly distribute all flows to available queues.

2 Contribution-aware Pausing



- ☐ Records the size of each flow in Flow Table, indexed by *hash(FID)*
- ☐ Pausing Decision:
 - \square If $Q_h > Q > Q_l$: flow with size > fair size should be paused.
 - \square If $Q > Q_h$, all passed flow should be paused.
- \square Fair size: $[L >> [log_2N]]$.

Opportunistic Re-assigning



- PAUSE carries FID
- \square Re-direct all congested flows to a reserved isolation queue (rsvQ) by controlling the flow-to-queue mapping in \blacksquare .
- \square Resuming condition of rsvQ:
 - ☐ 1) all isolated flows have received RESUME;
 - ☐ 2) previous buffered packets have been drained off

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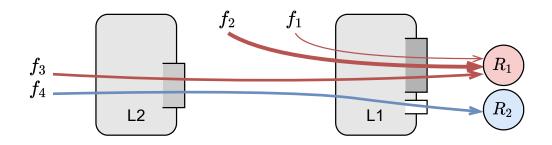
Micro-benchmark

Setting

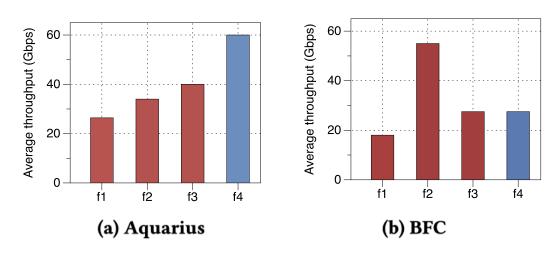
- NS-3 simulator
- f1: 33Gbps; f2~ f4: 100Gbps.
- R1 becomes the bottleneck.

Aquarius achieves perflow control granularity.

- Fair partition of bottleneck link capacity between f1 to f3.
- The victim flow, f4, is not affected.



Micro-benchmark setting.

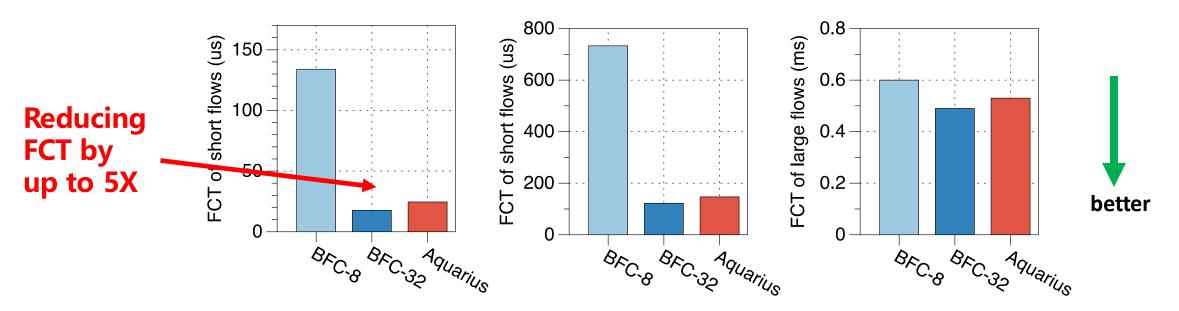


Average throughput for flows $f1 \sim f4$.

Realistic Traffic

Setting

- NS-3 simulator; 3-layer fat-tree topology; 48 switches; 128 servers
- 100Gbps link; 1us propagation delay; 12MB switch buffer
- Web Server with a 70% average load and 5% 100-to-1 incast traffic



FCT under Web Server distribution with 70% load and 5% 100-1 incast.

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Summary

■ Per-hop Flow Control is Necessary but Prior Scheme is Insufficient

- End-to-end CC alone is insufficient for managing transient congestion.
- Prior flow control schemes either lack fine-grained control or require an impractical number of queues.
- BFC experiences considerable performance degradation when queues are limited.

Key Idea of Aquarius

To approximate the ideal flow control behavior without requiring per-flow queues.

Key points of Aquarius

- Contribution-aware pausing, that accurately identifies the set of congested flow, mimics the behavior of ideal FC at the congested port.
- Opportunistic Re-assigning, that isolates congested flows from normal queues, mimics the behavior of ideal FC at the upstream port.

Thank you!

Contact email: wlicv@connect.ust.hk

Order Mark

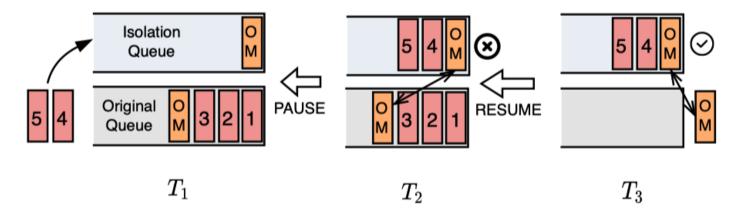


Fig. 7: A high-level view of how Order Mark (OM) packets support in-order delivery.